# A domain-wall encoding of discrete variables* 

## Zaiku Group

Nicholas Chancellor

March 12, 2021



[^0]
## One-hot constraints

- Common constraints found in optimisation problems
- Enforce that exactly one mutually exclusive option is taken
- Examples: a vehicle can only take one route to get to a destination, a piece of equipment can only be doing one thing at a time, etc...
Easy to enforce using quadratic interactions, term proportional to $\left(\sum_{i} x_{i}-1\right)^{2}$ equal to zero if exactly one variable is 1 and all others are 0 , higher otherwise
...but requires interaction between all variables



## A slightly different perspective: discrete variables

- Rather than thinking of individual binary variables under a constraint, treat like a single $m$ value variable
- A mostly philosophical distinction, but crucial for understanding other encodings
- Define two index objects:

$$
x_{i, \alpha}= \begin{cases}1 & \text { variable } d_{i} \text { takes value } \alpha \\ 0 & \text { otherwise }\end{cases}
$$

- Discrete Quadratic models, (DQM), made from pairwise interactions of $x$ terms:

$$
H_{\mathrm{DQM}}=\sum_{i, j} \sum_{\alpha, \beta} D_{(i, j, \alpha, \beta)} x_{i, \alpha} x_{j, \beta}
$$

## one-hot and binary encoding as a DQM

$$
H_{\text {one hot }}=H_{\mathrm{DQM}}+\lambda \sum_{i}\left(\sum_{\alpha=0}^{m-1} x_{i, \alpha}-1\right)^{2}
$$

- Each $x_{i, \alpha}$ is an individual binary variable, add one-hot constraints to force them to be single valued
- Easy to forget distinction in this case, but in general $x_{i, \alpha}$ does not need to map to a single binary variable
- Example: binary encoding, each binary variable $b \in\{0,1\}$ is a digit in a binary number $x_{i, \alpha}$ is a string of $\log _{2}(m) b$ and $1-b$ terms


## Is there another way?

YES!

1. Constrain Ising $s_{i, \alpha} \in\{-1,+1\}$ variables with strong ferromagnetic coupling

$$
H_{\mathrm{chain}}=-\kappa\left(\sum_{\alpha=-1}^{m-2} s_{i, \alpha} s_{i, \alpha+1}\right)
$$

2. Constrain $s_{i,-1}=-1$ and $s_{i, m-1}=1$ so that the chain consisting of variables $1 \ldots m-2$ is frustrated and contains at least one domain wall
3. Define DQM terms


## Domain walls used to store information

Simple example with four Ising variables:

| encoded value | qubit configuration |
| :---: | :---: |
| 0 | 1111 |
| 1 | -1111 |
| 2 | $-1-111$ |
| 3 | $-1-1-11$ |
| 4 | $-1-1-1-1$ |



## Aside: what if I encode a binary variable using domain walls?

Recover normal binary encoding, why?

- Take $m-1$ qubit segment of chain, that is only a single qubit
- The coupling to "virtual" qubits on either side cancels out
- DQM terms recover Ising definitions for 2-SAT clauses

Domain-wall encoded DQM variables can be used in problems which also contain normal binary variables*
*The synthetic problems in Chancellor, Phys. Rev. A 102, 062606 are an example of this

How does this domain wall encoding stack up against one-hot?

- One fewer qubit per DQM variable than one-hot
- $x_{i, \alpha}$ is linear in $s$ terms, therefore $x_{i, \alpha} x_{j, \beta}$ is quadratic, maps to Ising model with only second order interactions
- Simple degree of freedom counting arguments $\rightarrow$ domain-wall encoding uses the smallest possible number of variables for all $x_{i, \alpha} x_{j, \beta}$ terms to be quadratic
- Less connectivity within variables, linear versus all-to-all




## Key stats for implementation

Note $d_{e}=$ edge distance, minimum number of interaction edges to cross to get from one qubit to another

| performance metric | binary | one-hot | domain wall |
| :---: | :---: | :---: | :---: |
| \# qubits | $\left\lceil\log _{2}(m)\right\rceil$ | $m$ | $m-1$ |
| \# couplers <br> for encoding | 0 if $m=2^{n} n \in \mathbb{Z}$ <br> complicated otherwise | $m(m-1)$ | $m-2$ |
| intra-variable connectivity | $\mathrm{N} / \mathrm{A}$ or complicated | complete | linear |
| maximum order <br> needed for two variable interactions | $2\left\lceil\log _{2}(m)\right\rceil$ | 2 | 2 |
| maximum $d_{e}$ between <br> qubits in interacting variables | complicated | 2 | $m$ |

Last row, domain wall variables allowed to be more "spread out" may be easier to minor embed

## Test this on an example: colouring problems^

Simple test problem with structure: penalty between nodes if and only if they are the same colour
Use natural structure of problem to 'spread out' embedding
Four colouring example, 'layered' structure in Domain wall (right), no structure in one hot, (left)

three-colouring $\rightarrow$ randomly generated edges with $50 \%$ probability k -colouring $\rightarrow$ twice as many nodes as colours, random edges with $75 \%$ probability

[^1]
## Test this numerically with minor-miner*



Embedding ratio $=$ number of qubits per logical variable ( Y -axis=one-hot, X -axis=domain-wall) chimera in red, Pegasus in blue
(a) Max-three-colouring problems (domain-wall slightly worse)
(b) Max-k-colouring problems (domain-wall much better)
(c) Artificial scheduling problem (domain-wall much better)
*Example taken from Chancellor, Quantum Sci. Technol. A 045004

Bigger problems can be embedded on the same size device*


Size of device required to embed problems of different sizes right column $=$ three colouring, left= $k$-colouring

- Can embed larger problems for both graphs and problem types
- Fewer binary variables makes up for slightly worse embedding ratio in three-colouring case

[^2]
## What about dynamics?

- Makes embedding more efficient...
- But does it translate to performance gains on actual annealers? how does the encoding affect the dynamics?

Not easy to answer a priori, on one hand

- Changing the values of domain-wall variables is non-perturbative*
- This isn't true for one-hot

On the other hand

- Configurations now have a defined order and are traversed in a linear way, rather than one-hot where there is no "order"
- Not clear if this will be helpful or harmful

Need to test experimentally, our recent paper ar $\chi$ iv:2102.12224 does this

[^3]
## The results^

For both k and three colouring problems the domain-wall encoding performs better on both Advantage and 2000Q
three colouring (left), k-colouring (right)


$\mathrm{C}=$ number of places same colour touches
Even looks like domain-wall on 2000Q out-performs one-hot on Advantage!
Use hypothesis testing to verify that this is a statistically significant result, test 100 instances on each and see how much each processor/encoding combination wins for all 6 combinations

[^4]
## Hypothesis testing, three colour*

Green=statistically significant result (95\% confidence)


- Domain wall 2000Q beats one hot Advantage (in a statistically significant way)
- Trend continue up to size where no longer possible to embed in 2000Q (FAIL)
- Otherwise results are expected $\rightarrow 2000$ Q worse than Advantage, one hot worse than domain wall

[^5]
## Hypothesis testing, k colour*

Green/red=statistically significant result (95\% confidence)

|  | Adv. dw/oh |  | 2000Q dw/oh |  | dw Adv./2000Q |  | oh Adv./2000Q |  | (dw, Adv.)/(oh, 2000Q) |  | (dw, 2000Q)/(oh, Adv.) |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 3 color (b,w) | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 3 color p |  |  |  |  |  |  |  |  |  |  |  |  |
| 4 color (b,w) | 34 | 1 | 37 | 2 | 11 | 3 | 26 | 16 | 44 | 1 | 33 | 7 |
| 4 color p | $1.05 \times 10^{-9}$ |  | $1.42 \times 10^{-9}$ |  | $2.87 \times 10^{-2}$ |  | $8.21 \times 10^{-2}$ |  | $1.31 \times 10^{-12}$ |  | $2.11 \times 10^{-5}$ |  |
| 5 color (b,w) | 91 | 1 | 78 | 1 | 34 | 18 | 23 | 59 | 88 | 1 | 91 | 1 |
| 5 color p | $1.88 \times 10^{-26}$ |  | $1.32 \times 10^{-22}$ |  | $1.82 \times 10^{-2}$ |  | $\approx 1$ |  | $1.45 \times 10^{-25}$ |  | $1.88 \times 10^{-26}$ |  |
| 6 color(b,w) | 99 | 0 |  | FAIL | 59 | 15 |  | FAIL |  | FAIL | 99 | 0 |
| 6 color p | $1.58 \times 10^{-30}$ |  | $1.28 \times 10^{-7}$ |  |  |  |  |  |  |  |  |  |
| 7 color(b,w) | 92 | 0 | FAIL | FAIL |  | FAIL |  | FAIL |  | FAIL | FAIL |  |
| 7 color p | $2.02 \times 10^{-28}$ |  |  |  |  |  |  |  |  |  |  |  |

- Domain wall 2000Q beats one hot Advantage (in a statistically significant way)
- Trend continue up to size where no longer possible to embed in 2000Q (FAIL)
- One case where 2000Q beats advantage for the same decoding (one hot) ${ }^{\star}$
*This goes away when the decoding strategy for broken chains is changed so probably an artefact of majority vote decoding
*arұiv:2102.12224


## Same pattern holds for probability to find optimal ${ }^{\star}$

three colouring (left), k-colouring (right)


Note that each run was only performed with 100 reads, better results could be attained with more reads

All QPU-encoding combinations found optimal solution at smallest size $\rightarrow$ explains no "winners" in hypothesis testing

## Digging deeper into performance: encoding failures ${ }^{\star}$

What fraction of solutions have all one-hot/domain-wall constraints satisfied
three colouring (left), k-colouring (right)



Domain-wall constraints are much less "fragile" especially with only three colours, makes a much bigger difference than processor structure

[^6]
## Digging deeper into performance: chain breaks^

What fraction of solutions have no unbroken minor embedding chains
three colouring (left), k-colouring (right)



Note: bars are standard deviation, standard error is $10 x$ smaller
QPU structure seems to make a bigger difference here, but domainwall encoding still leads to an improvement

[^7]
## Experimental results summary

- Encoding makes a bigger difference to solution optimality even than choosing a more advanced processor
- Domain wall constraints seem much less "fragile"
- Encoding still helps with chain breaks, but advantage is smaller $\rightarrow$ QPU structure makes a bigger difference


## Experiments didn't find any metrics where one-hot does better

No observed downside to using domain-wall encoding, but some major advantages

## Want to try it yourself?

Python code to create domain wall encodings available at https://collections.durham.ac.uk/: "Domain wall encoding of integer variables for quantum annealing and QAOA [dataset]"*


[^8]
## How to use the repository code

- Load the ez_domain_wall module
- The function make_domain_wall_encoding creates a domain wall encoding of a discrete problem
- ez_dw_examples jupyter notebook with some examples of how to use the code

Inputs to make_domain_wall_encoding

1. Variable sizes: a list of the sizes of all variables, $\rightarrow$ call the sum of these sizes $R$
2. Penalties: single body terms which penalize different values of individual variables, 1-D array (or list) of length $R$
3. Interactions: interactions between variables upper triangular $R$ by $R$ array

Outputs: J_core, h_core terms for enforcing domain wall constraints, J_prob, h_prob rest of Hamiltonian

## A simple example: three colouring a line of three nodes

1. Three variables with three possible colours, therefore variable sizes $\rightarrow[3,3,3](R=3 \times 3=9)$
2. No single term penalties, therefore penalties term is all zeros (length 9)
3. Interaction terms: $9 \times 9$ array of zeros with 6 terms which take value 1 (using Python style zero indexing) $(0,3),(1,4),(2,5) \rightarrow$ prevent the first and second node from being the same colour
$(3,6),(4,7),(5,8) \rightarrow$ prevent the second and third node from being the same colour

## The outputs for our simple example

$$
\begin{aligned}
& \text { h_core: } {\left[\begin{array}{cccccc}
1 & -1 & 1 & -1 & 1 & -1
\end{array}\right] } \\
& \text { J_core: } {\left[\begin{array}{cccccc}
0 & -1 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & -1 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & -1 \\
0 & 0 & 0 & 0 & 0 & 0
\end{array}\right] } \\
& \text { h_prob: } {\left[\begin{array}{llllll}
0.25 & -0.25 & 0.5 & -0.5 & 0.25 & -0.25
\end{array}\right] } \\
& \text { J_prob: } {\left[\begin{array}{cccccc}
0.0 & 0.0 & 0.5 & -0.25 & 0.0 & 0.0 \\
0.0 & 0.0 & -0.25 & 0.5 & 0.0 & 0.0 \\
0.0 & 0.0 & 0.0 & 0.0 & 0.5 & -0.25 \\
0.0 & 0.0 & 0.0 & 0.0 & -0.25 & 0.5 \\
0.0 & 0.0 & 0.0 & 0.0 & 0.0 & 0.0 \\
0.0 & 0.0 & 0.0 & 0.0 & 0.0 & 0.0
\end{array}\right] } \\
& H=\sum_{i<j} J \_p r o b_{i j} s_{i} s_{j}+\sum_{i} h_{-} p r o b_{i} s_{i}+ \\
& \kappa\left(\sum_{i<j} J_{-} \operatorname{core}_{i j} s_{i} s_{j}+\sum_{i} h_{-} c o r e_{i} s_{i}\right) \\
& s \in\{-1,1\}
\end{aligned}
$$

## Possible application: quantum simulation

- Treat each variable in the DQM as a point in space
- Quadratic coupling can emulate many differential equation terms (may also be useful in solving other diff. eqs)
- Transverse fields can emulate local quantum fluctuations in quantum field theories (field fluctuates to nearby values)

Under development with Steven Abel and Michael Spannowsky in Durham IPPP*
Not the main focus of the talk, but worth highlighting

[^9]
## Further outlook (Chancellor, QST. 4 045004)

Drivers which preserve valid subspace out of two body terms ${ }^{\star}$

- Constructed by combining $Z_{i-1} X_{i}$ and $-X_{i} Z_{i+1}$ terms $\rightarrow$ rotations only happen if domain wall is present
- May be useful in QAOA or other gate model algorithm (maybe even just with transverse field drivers)
- Annealers with muti-body drivers (longer term)

Layered structure for important problem types such as colouring and scheduling

- Suggests application specific (ASIC if superconducting) designs for future annealers


[^10]
## Summary

Domain-wall encoding can lead to superior performance over one-hot for quantum annealing

- No downside seen yet
- Can get substantial gains, more than new hardware in the cases we tested

Should try it if you are using discrete variables

- Python code available so you don't have to code "from scratch"
- dw encoding to be included in DLR (German aerospace organization) software package when released
- QAOA has not been tested, not clear how it would perform


## Potential for using annealers as simulators

- Interesting early results, still being explored
- More rigorous understanding may be helpful


[^0]:    *Based on results from ar $\chi$ iv:2102.12224 (with co-authors Jie Chen and Tobias Stollenwerk) and background from other sources

[^1]:    *Example taken from Chancellor, Quantum Sci. Technol. 4 045004

[^2]:    *Example taken from Chancellor, Quantum Sci. Technol. 4 045004

[^3]:    *all valid configurations can be reached without having to pass through invalid configurations

[^4]:    *arұiv:2102.12224

[^5]:    *ar $\chi$ iv:2102.12224

[^6]:    *arұiv:2102.12224

[^7]:    *arұiv:2102.12224

[^8]:    *https://doi.org/10.15128/r27d278t029

[^9]:    *see: S. Abel, N. Chancellor, M. Spannowsky Phys. Rev. D 103, 016008 (2021) and S. Abel, M. Spannowsky ar $\chi$ iv:2006.06003

[^10]:    *analogous to what was done in Hadfield et. al. Algorithms 12.2 (2019):34

